
Security Evaluation of Biometric Privacy Enhancing Techniques

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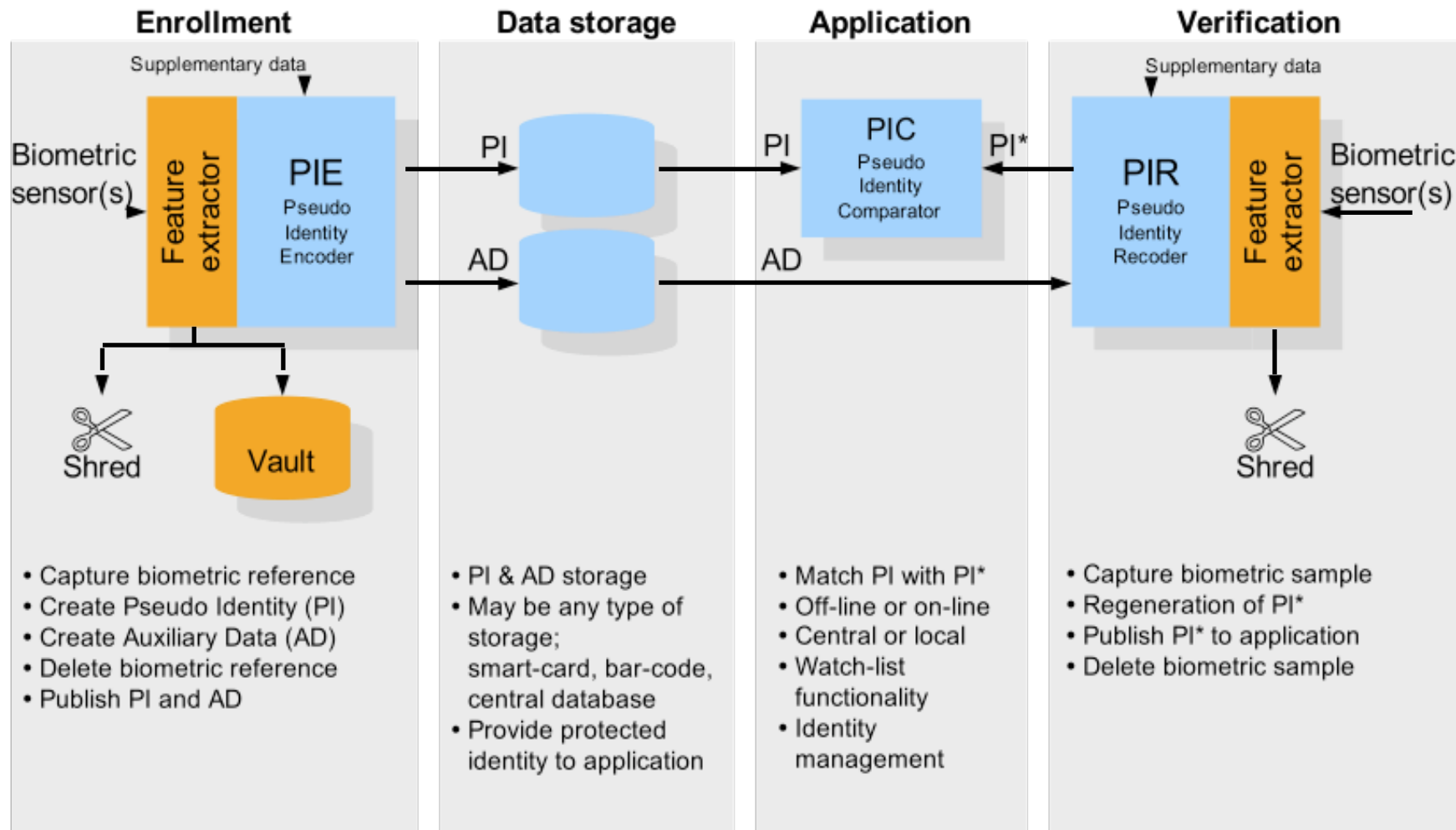
Motivation

- Privacy and Security Risks in Common Biometric System
 - Identity theft
 - Cross matching
 - Unchangeability
 - Privacy and legislation
- Template protection technique
 - It converts biometric data into **multiple various** references, from which it is **infeasible and hard** to retrieve the original information.
 - Biometric template protection is a very important supplement to improve security and enhance privacy protection of biometric systems.
- New challenge: how to evaluate *security* and *privacy enhancing ability* of a template protection technique?

• ISO/IEC CD 24745 Information technology - Security techniques - Biometric template protection



Reference Architecture for Template Protection Technique



• J. Breebaart, C. Busch, J. Grave, E. Kindt: A Reference Architecture for Biometric Template Protection based on Pseudo Identities, in Proceedings BIOSIG 2008



Template Protection Techniques*

- Protected template
 - Pseudonymous Identifier : PI
 - Auxiliary Data: AD
- Supplementary Data: SD
- Pseudonymous Identifier Encoder:
 - $[PI, AD] = PIE(X, SD)$
- Pseudonymous Identifier Recorder:
 - $[PI'] = PIR(X', AD, SD)$



State of the Art of Template Protection Techniques

■ Transformation algorithms

■ Biometric Salting:

- *Biobhashing*: randomly converts features into bit strings

- *Biometric Encryption*: whitening features in frequency space

- Cancelable Biometrics: transformation of features or morphing of samples with non-invertible function

■ Biocrypto algorithms

- *Fuzzy commitment*: cryptographic hashing (ordered set)

- *Fuzzy vault*: secret sharing protocol (non-ordered set like fingerprint minutiae)

State of the Art of Template Protection Techniques

		Transformation	Biocrypto
Protected template (public)	PI	Transformed feature/sample	Protected binary identifier
	AD	--	e.g. helper data or point set
SD (secret)		Specific parameters used in transformation function	Optional personal PW
Comparison result		Similarity score / Exact match	Exact match

Protection Goals

- Sufficient security of PI:
 - It is hard to guess the true binary identifier for an enrolled subject in the biocrypto system
 - It is hard to reconstruct a sample or feature set yielding an identical transformation, which can successfully pass pseudonymous identifier verification
- Low leakage of biometric information in PI or AD
 - Little biometric information can be learned from PI or AD
- Good randomness of PI
 - The true binary identifiers should be independent
 - No correlation between PIs (the transformed samples/features)
- No personally identifiable information in AD
 - AD is at best random



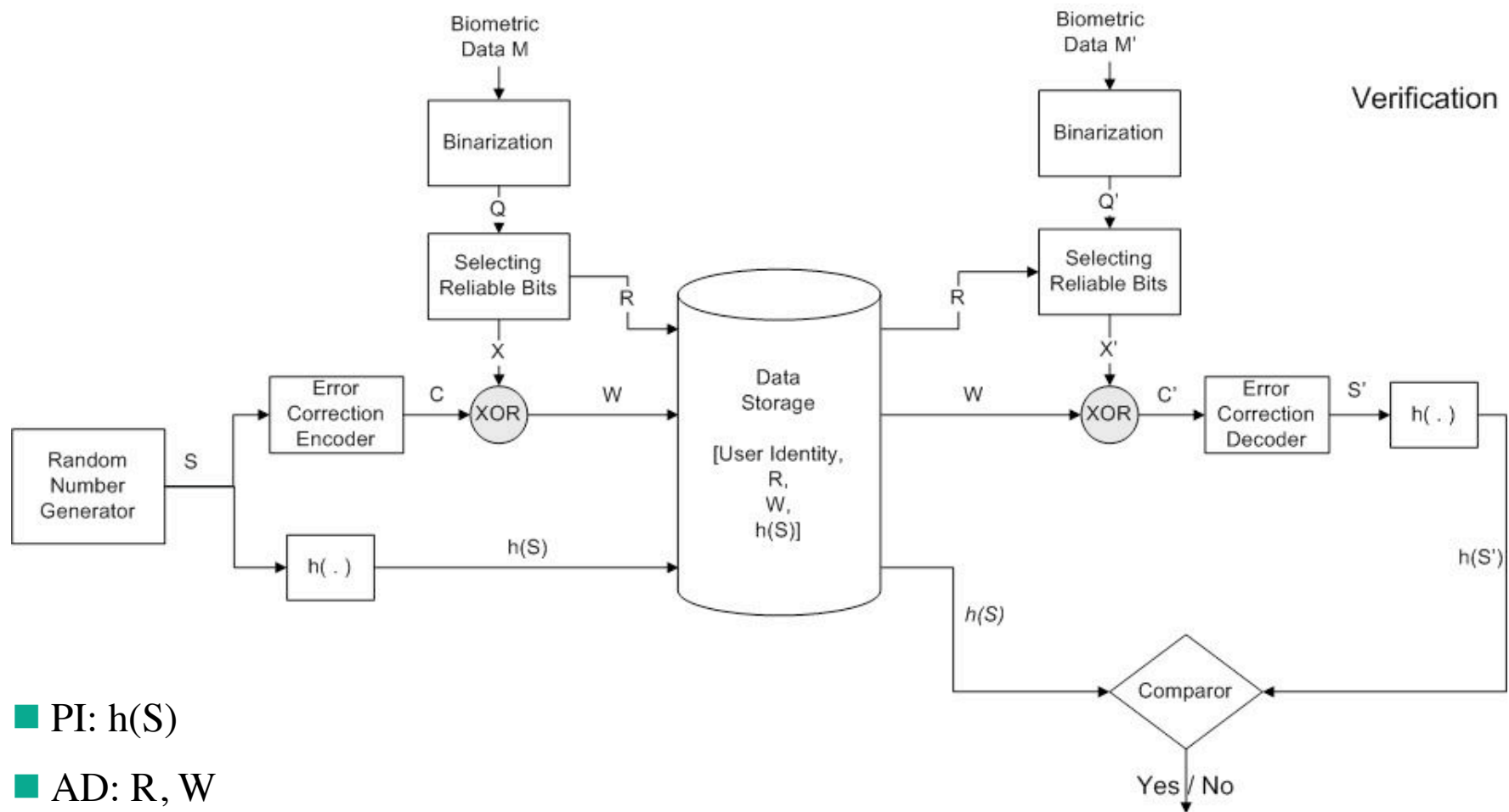
Threat Models (TM)

- TM 1: Storage - A protected template of a subject is known

- TM 2: Signal processing - Template protection algorithm is public (Kerckhoffs' principle)
 - Linkage between databases - Stored protected templates of the same subject from different applications are known
 - System parameters are known
 - (Secret) supplementary data SD are known
 - The statistical properties of the biometric data are known

- TM 3: Decision - A large biometric database is available

An Example of Fuzzy Commitment Scheme



■ PI: $h(S)$

■ AD: R, W

■ System information: coding algorithm, coding parameters, e.g. L_s, L_c

Security Evaluation of the Fuzzy Commitment Scheme

- TM 1: A protected template is known
 - Cryptographic security of $PI = L_s$ bits (secret size)
 - Leakage of biometric information in PI and AD is neglectable
- TM 2: Template protection algorithm is public
 - Coding algorithm and coding parameters are known:
 - Cryptographic security of $PI = L_s$ bits
 - Leakage of biometric information in $AD = L_c - L_s$ bits
 - More entries of the same subject in different databases are known:
 - Identification ability of AD : FMR, FNMR or probability of identification for numerous ranks given R (R is indicating reliable bits in the binary vector)

Security Evaluation of the Fuzzy Commitment Scheme

- TM 2 (continued): Template protection algorithm is public
 - Coding algorithm, parameters and more entries of the same subject in different databases (combining leakage)
 - Cryptographic security of PIs: $\hat{H}(S1|W1, W2), \hat{H}(S2|W1, W2)$
 - Security leakage: $L_s - \hat{H}(S1|W1, W2), L_s - \hat{H}(S2|W1, W2)$
 - Leakage of biometric information in AD: $\hat{I}(X1; W1, W2), \hat{I}(X1; W1, W2)$
 - Statistical properties of binary biometric feature vector are known or partially known as well as the coding algorithm and coding parameters:
 - Cryptographic security of PI: $\hat{H}(S|W)$
 - Security leakage: $L_s - \hat{H}(S|W)$
 - Leakage of biometric information in AD: $\hat{I}(X; W)$

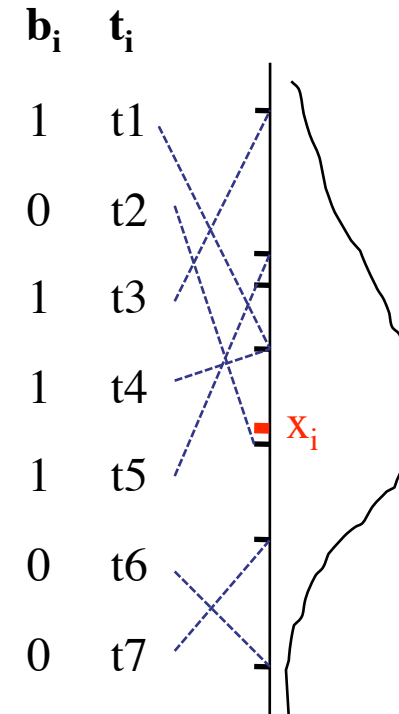
Security Evaluation of the Fuzzy Commitment Scheme

■ Remarks:

- A stream cipher is used in the fuzzy commitment scheme
 - Cryptographic security of PI = Uncertainty about X in W ($H(X|W)$)
 - The assumption that the public helper data W has zero leakage of S is true only in the case that X is identically independently distributed
 - The security leakage is allowed, if PI is cryptographically secure
 - Leakage of biometric information in AD can lead to leakage of PI: $\hat{H}(S1|W1, W2) \leq \hat{H}(S1|W1)$. Because W2 is dependent on biometric information X2 and randomly generated S2, increasing leakage is caused by the correlation between W2 and X1

An Example of Biohashing based on Scalar Randomization

- Any feature x_i in $\mathbf{x} = \{x_i | i \in [1, \dots, n]\}$ can be binarized with m randomly generated thresholds $\mathbf{t} = \{t_i^j | j \in [1, \dots, m], i \in [1, \dots, n]\}$ into $\mathbf{b} = \{b_i^j | j \in [1, \dots, m], i \in [1, \dots, n]\}$, where t_i^j has the same distribution as x_i .
- Comparison based on Hamming distance between Biohashes.
- Pseudonymous Identifier PI : \mathbf{b} containing $m \times n$ bits
- Supplementary Data SD : \mathbf{t}
- System parameter: m, n



Security Evaluation of the Biohashing Algorithm

- TM 1: A protected template (Biohashes) is known
 - --: without information about Biohashing algorithm, Biohashes mean nothing

- TM 2: Template protection algorithm is public

- The parameters of the algorithm are known:

- Leakage of biometric information in PI: $\hat{f}(x_i) = \frac{\sum_j b_i^j}{m}$, where $\hat{f}(x_i)$ is the cumulative probability of with e.g. the 95% confident interval of

$$\hat{f}(x_i) \pm 1.96 \sqrt{\frac{\hat{f}(x_i)(1 - \hat{f}(x_i))}{m}}$$

- Randomness of PI = $\sum_{i=1}^n m \cdot H(f(x_i))$; for all users, the average information rate of Biohashes is 0.72



Security Evaluation of the Biohashing Algorithm

■ TM 2 (continued):

- The parameters of the algorithm and SD (binarization thresholds are compromised) are known

- Leakage of biometric information

- $$\frac{|\min_{b_i=1}\{t_i\} - \max_{b_i=0}\{t_i\}|}{|\max\{t_i\} - \min\{t_i\}|}$$

- $$\hat{p}(x_i) = \hat{p}(t_i)$$

- If more entries of the same subject in different databases are given, more accurate estimation can be done

- The parameters of the algorithm and the distribution of x are known

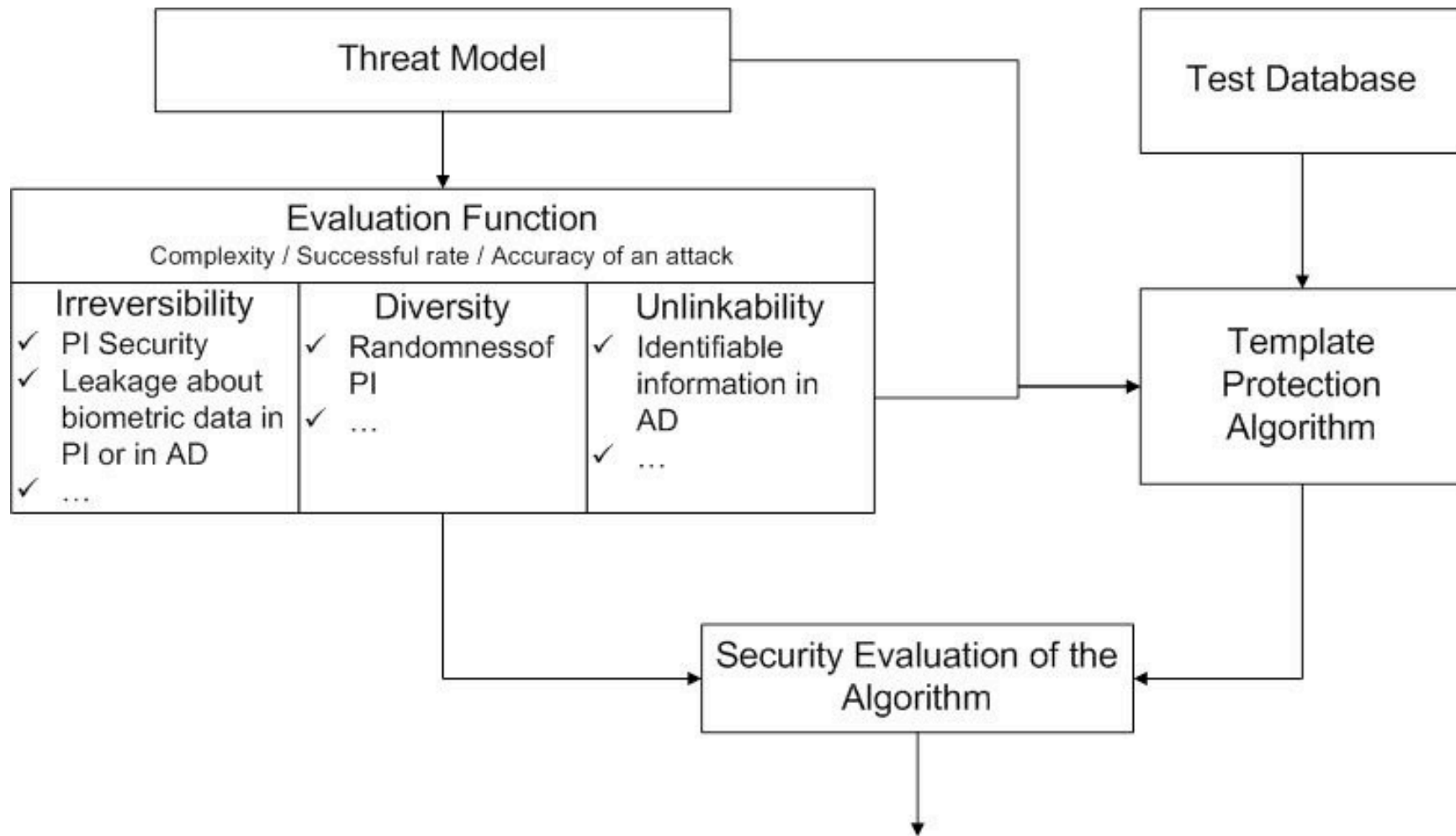
- Leakage of biometric information in PI:
$$\hat{x}_i = f^{-1}\left(\frac{\sum_j b_i^j}{m}\right)$$



Threat Model 3

- False acceptance attack: $P(PI_X = PIR(X', AD, SD)) = FMR$, where X and X' are from different subjects.
- FMR shows the average collision probability of biometric data
- The probability of successful false acceptance attack might be different than FMR. For example in Doddington's biometric menagerie:
 - Lambs: $FMR_L > FMR$
 - Wolves: $FMR_W < FMR$
 - In contrast to randomly generated identifiers, where all the identifiers are equally probable, pseudonymous identifiers generated from lambs/wolves characteristics are not equally probable
- False acceptance is inherent prosperity of biometric system and false acceptance attack is not avoidable. However it can be controlled with e.g. time out policy

Evaluation Framework



Remarks on Evaluation Framework

- All the requirements on template protection should be measured quantitatively with evaluation function
- From the security and privacy point of view, leakage on biometric information should be limited
 - to AD of protected template in biocrypto algorithms
 - even when the SD of transformation algorithms is compromised
- The likelihood to reconstruct true biometric data from protected template is dependent on SD chosen. An efficient template protection algorithm should have homogenous security for all the subjects and possible SD .

Conclusions

- A general evaluation framework according to 3 main threat models was developed.
- Depending on the threat model, the **measurable** security evaluation functions are defined
- It is necessary to measure privacy leakage of protected template including leakage of biometric information and any other personal identifiable information. Privacy leakage causes linkability and can give rise to further security problem
- In order to improve confidentiality of authentication and integrity of protected templates, cryptographic technique such as digital signature and encryption can be combined with template protection algorithms.

Thank you for your attention

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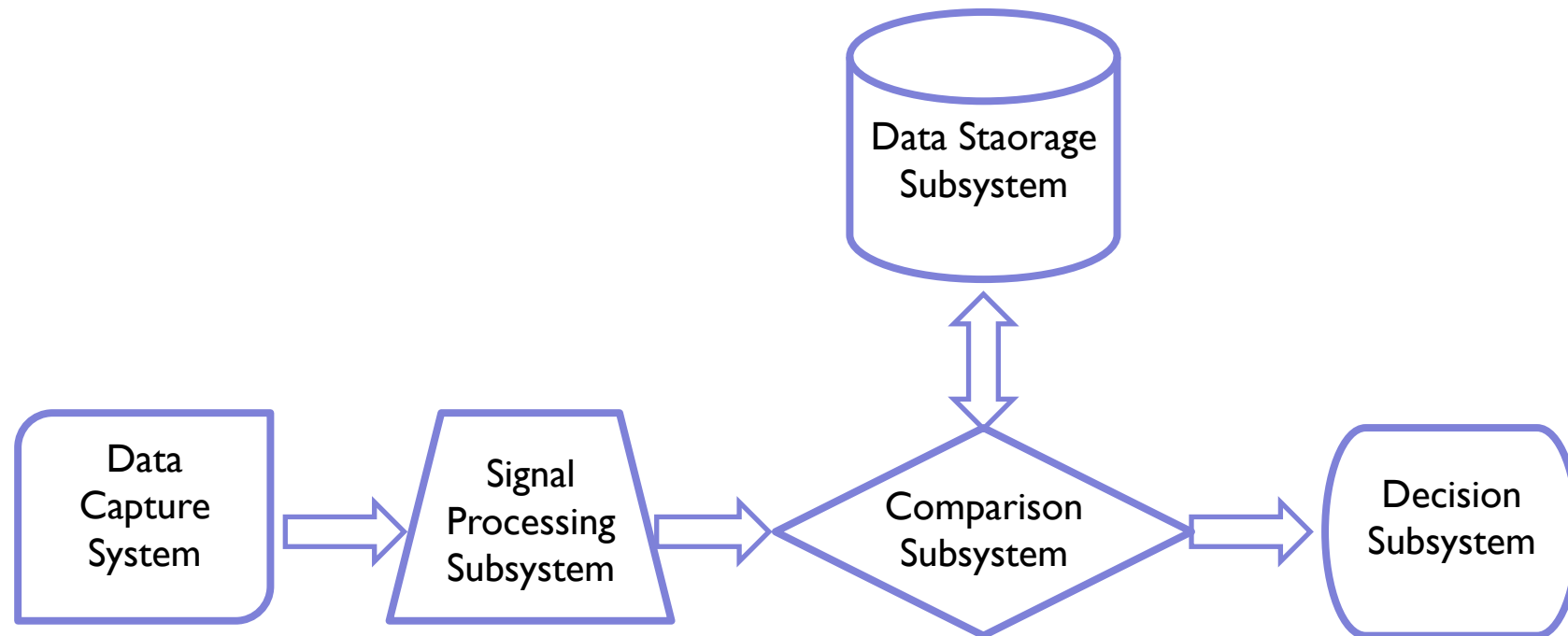
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Threats in the Biometric System



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